Sorting Algorithms

Chapter 9

CS 3358
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What is sorting?

- Sort: rearrange the items in a list into ascending or descending order
  - numerical order
  - alphabetical order
  - etc.

55 112 78 14 20 179 42 67 190 7 101 1 122 170 8

1 7 8 14 20 42 55 67 78 101 112 122 170 179 190

Why is sorting important?

- Searching in a sorted list is much easier than searching in an unsorted list.
- Especially for people
  - dictionary entries
  - phone book
  - card catalog in library
  - bank statement: transactions in date order
- Most of the data displayed by computers is sorted.

Sorting

- Sorting is one of the most intensively studied operations in computer science
- There are many different sorting algorithms
- The run-time analyses of each algorithm are well-known.

Sorting algorithms

- Selection sort
- Insertion sort
- Bubble sort

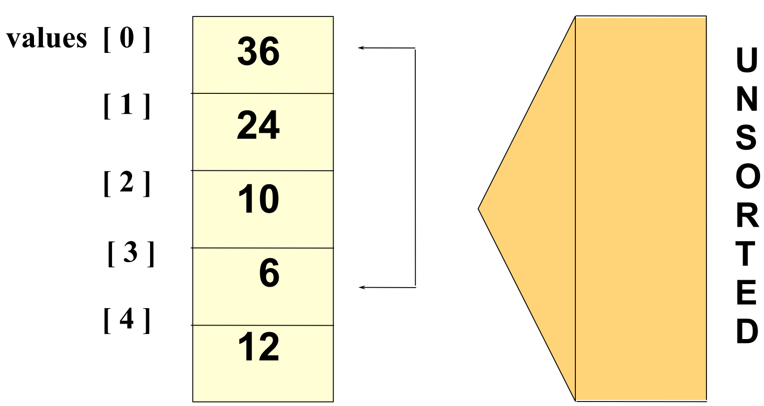
- Merge sort
- Quicksort

- Heap sort (later, when we talk about heaps)

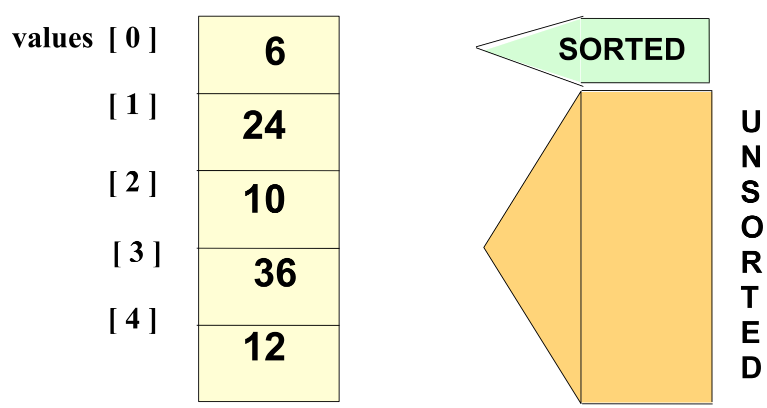
Selection sort

- There is a pass for each position (0..size-1)
- On each pass, the smallest (minimum) element in the rest of the list is exchanged (swapped) with element at the current position.
- The first part of the list (already processed) is always sorted
- Each pass increases the size of the sorted portion.

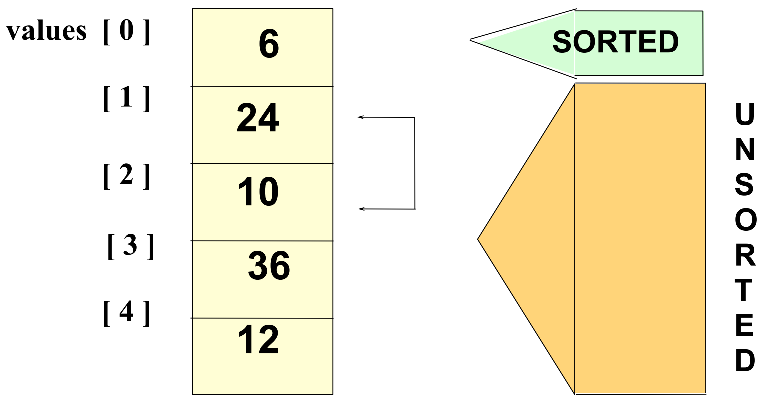
Selection Sort: Pass One



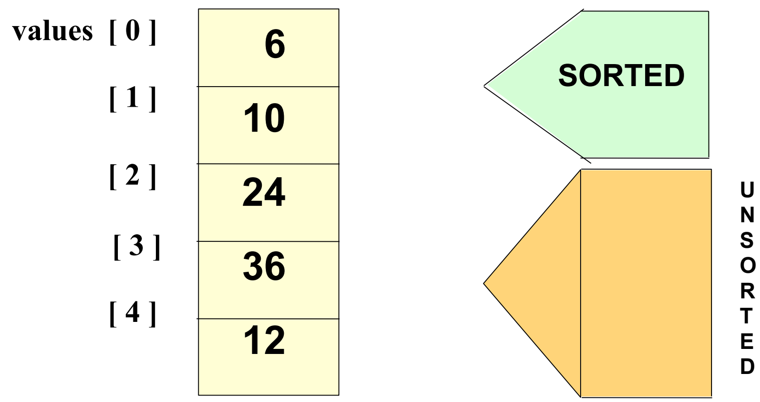
Selection Sort: End Pass One



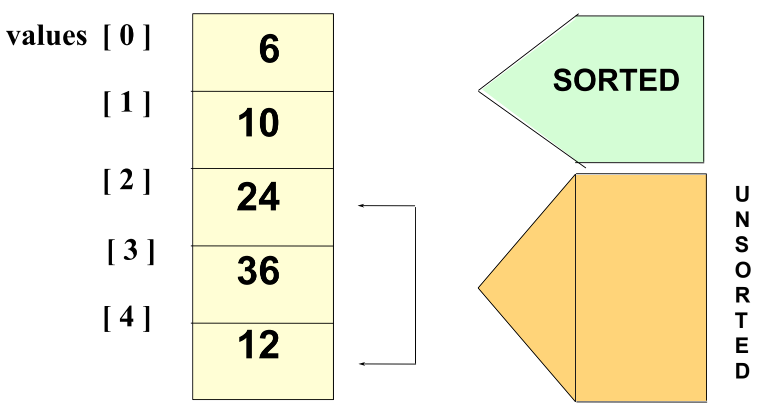
Selection Sort: Pass Two



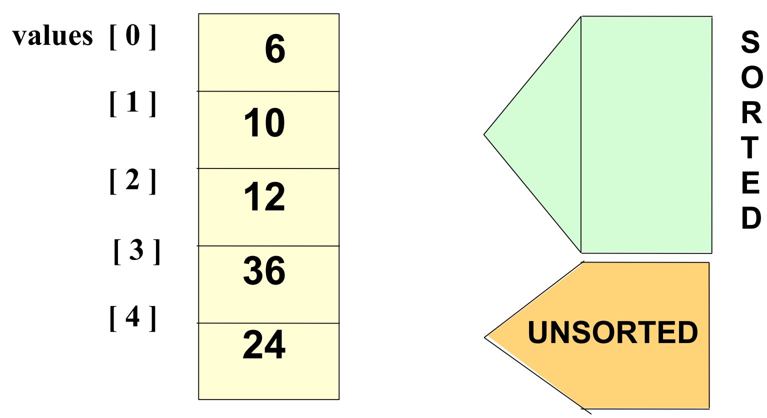
Selection Sort: End Pass Two



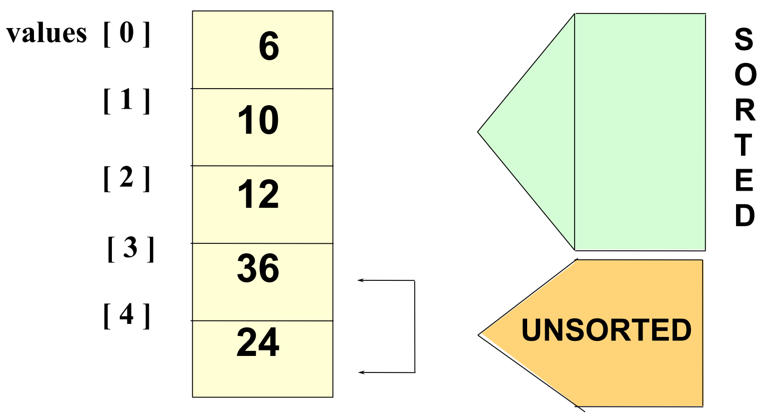
Selection Sort: Pass Three



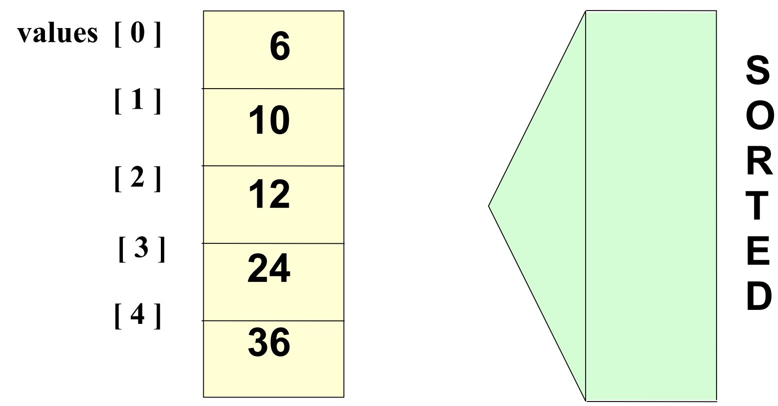
Selection Sort: End Pass Three



Selection Sort: Pass Four



Selection Sort: End Pass Four



Selection sort: code

```
template<class ItemType>
int minIndex(ItemType values[], int size, int start) {
    int minIndex = start;
    for (int i = start+1; i < size; i++)
        if (values[i] < values[minIndex])
            minIndex = i;
    return minIndex;
}

template<class ItemType>
void selectionSort (ItemType values[], int size) {
    int min;
    for (int index = 0; index < (size -1); index++) {
        min = minIndex(values, SIZE, index);
        swap(values[min],values[index]);
    }
}
```


Selection sort: runtime analysis

- N is the number of elements in the list
- Outer loop (in selectionSort) executes N times
- Inner loop (in minIndex) executes N-1, then N-2, then N-3, ... then once.
- Total number of comparisons (in inner loop):

(N-1) + (N-2) + . . . + 2 + 1

Note: N + (N-1) + (N-2) + . . . + 2 + 1 == N(N+1)/2

$=N^2/2 + N/2$

Subtract N from both sides: (N-1) + (N-2) + . . . + 2 + 1 $=N^2/2 + N/2 - N$

$=N^2/2 - N/2$

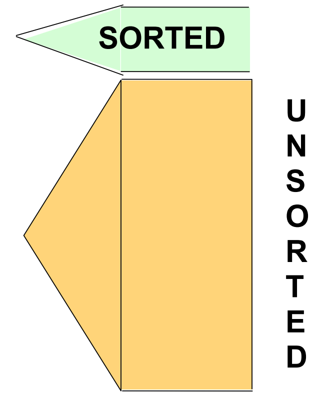
$\boxed{O(N^2)}$

Insertion sort

- There is a pass for each position (0..size-1)
- The front of the list remains sorted.
- On each pass, the next element is placed in its proper place among the already sorted elements.
- Like playing a card game, if you keep your hand sorted, when you draw a new card, you put it in the proper place in your hand.
- Each pass increases the size of the sorted portion.

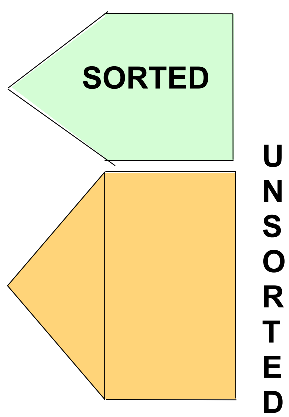
Insertion Sort: Pass One

values [0]	36
[1]	24
[2]	10
[3]	6
[4]	12



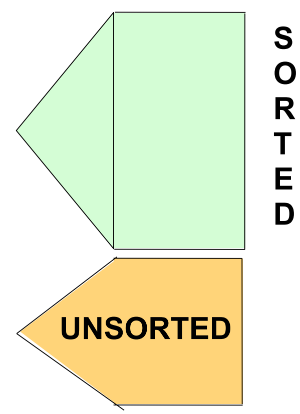
Insertion Sort: Pass Two

values [0]	24
[1]	36
[2]	10
[3]	6
[4]	12



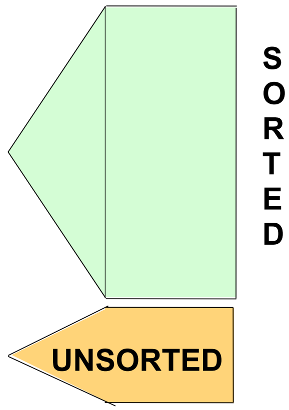
Insertion Sort: Pass Three

values [0]	10
[1]	24
[2]	36
[3]	6
[4]	12



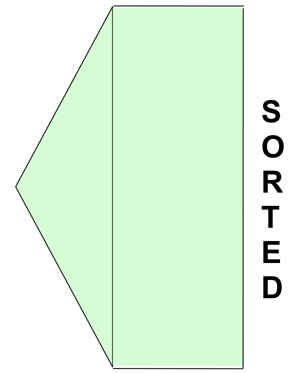
Insertion Sort: Pass Four

values [0]	6
[1]	10
[2]	24
[3]	36
[4]	12



Insertion Sort: Pass Five

values [0]	6
[1]	10
[2]	12
[3]	24
[4]	36



Insertion sort: code

```
template<class ItemType>
void insertionSort (ItemType a[], int size) {

    for (int index = 0; index < size; index++) {
        ItemType tmp = a[index];    // next element

        int j = index;        // start from the end

        // find tmp's place, AND shift bigger elements up
        while (j > 0 && tmp < a[j-1]) {
            a[j] = a[j-1];    // shift
            j--;
        }
        a[j] = tmp;           // put tmp in its place
    }
}
```


Insertion sort: runtime analysis

- Very similar to Selection sort
- Total number of comparisons (in inner loop):
  - At most j+1, which is 1, then 2, then 3 ... up to N
- So it's N + (N-1) + (N-2) + . . . + 2 + 1 == N(N+1)/2

$\boxed{O(N^2)}$

Bubble sort

- On each pass:
  - Compare first two elements. If the first is bigger, they exchange places (swap).
  - Compare second and third elements. If second is bigger, exchange them.
  - Repeat until last two elements of the list are compared.
- Repeat this process until a pass completes with no exchanges

Bubble sort
how does it work?

- At the end of the first pass, the largest element is moved to the end (it's bigger than all its neighbors)
- At the end of the second pass, the second largest element is moved to just before the last element.
- The back end (tail) of the list remains sorted.
- Each pass increases the size of the sorted portion.
- No exchanges implies each element is smaller than its next neighbor (so the list is sorted).

Bubble sort
Example

- 7 2 3 8 9 1   7 > 2, swap
- 2 7 3 8 9 1   7 > 3, swap
- 2 3 7 8 9 1   !(7 > 8), no swap
- 2 3 7 8 9 1   !(8 > 9), no swap
- 2 3 7 8 9 1   9 > 1, swap
- 2 3 7 8 1 <u>9</u>   finished pass 1, did 3 swaps

Note: largest element is in last position

Bubble sort
Example

- 2 3 7 8 1 9   2<3<7<8, no swap, !(8<1), swap
- 2 3 7 1 <u>8 9</u>   (8<9) no swap
- finished pass 2, did one swap

2 largest elements in last 2 positions

- 2 3 7 1 8 9   2<3<7, no swap, !(7<1), swap
- 2 3 1 <u>7 8 9</u>   7<8<9, no swap
- finished pass 3, did one swap

3 largest elements in last 3 positions

Bubble sort
Example

- 2 3 1 7 8 9   2<3, !(3<1) swap, 3<7<8<9
- 2 1 3 7 8 9
- finished pass 4, did one swap
- 2 1 3 7 8 9   !(2<1) swap, 2<3<7<8<9
- 1 2 3 7 8 9
- finished pass 5, did one swap
- 1 2 3 7 8 9   1<2<3<7<8<9, no swaps
- finished pass 6, no swaps, list is sorted!

Bubble sort: code

```
template<class ItemType>
void bubbleSort (ItemType a[], int size) {

    bool swapped;
    do {
        swapped = false;
        for (int i = 0; i < (size-1); i++) {
            if (a[i] > a[i+1]) {
                swap(a[i],a[i+1]);
                swapped = true;
            }
        }
    } while (swapped);
}
```


Bubble sort: runtime analysis

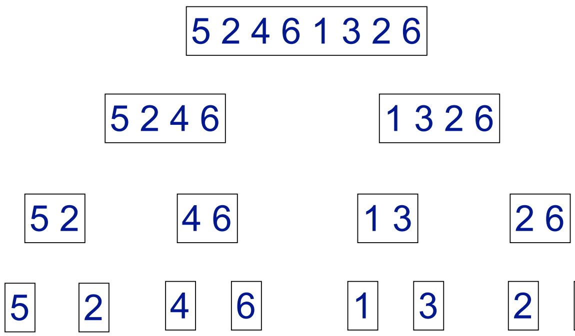
- Each pass makes N-1 comparisons
- There will be at most N passes
  - one to move the right element into each position
- So worst case it's: (N-1)*N $\boxed{O(N^2)}$
- What is the best case?
- Are there any sorting algorithms better than $O(N^2)$?

Merge sort

- Divide and conquer!
- 2 half-sized lists sorted recursively
- the algorithm:
  - if list size is 0 or 1, return (base case) otherwise:
  - recursively sort first half and then second half of list.
  - merge the two sorted halves into one sorted list.

Merge sort
Example

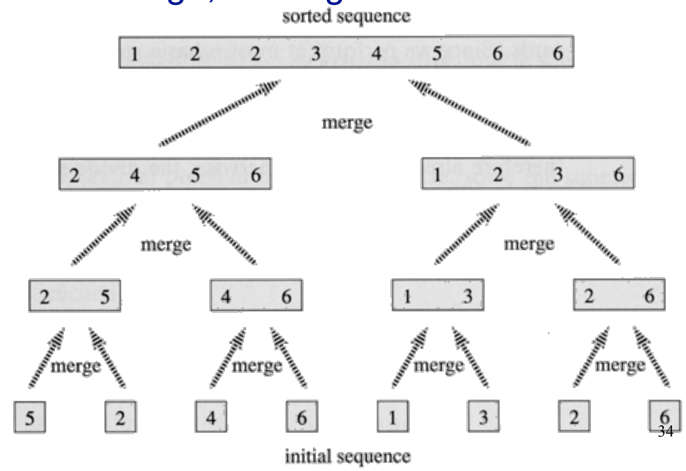
- Recursively divide list in half:



Each of these are sorted (base case length = 1)

Merge sort
Example

- Calls to merge, starting from the bottom:



Merge sort: code

```
template<class ItemType>
void mergeSort (ItemType values[], int first, int last) {
    if (first < last) {
        int middle = (first + last) / 2;

        mergeSort(values, first, middle);
        mergeSort(values, middle+1, last);

        merge(values, first, middle, last);
    }
}

template<class ItemType>
void mergeSort (ItemType values[], int size) {
    mergeSort(values, 0, size-1);
}
```


Merge sort: code: merge

```
template<class ItemType>
void merge(ItemType values[], int first, int middle, int last) {

    ItemType tmp[last-first+1];  //temporary array
    int i=first;        //index for left
    int j=middle+1;     //index for right
    int k=0;            //index for tmp

    while (i<=middle && j<=last)   //merge, compare next elem from each array
        if (values[i] < values[j])
            tmp[k++] = values[i++];
        else
            tmp[k++] = values[j++];

    while (i<=middle)            //merge remaining elements from left, if any
        tmp[k++] = values[i++];

    while (j<=last)              //merge remaining elements from right, if any
        tmp[k++] = values[j++];

    for (int i = first; i <=last; i++) //copy from tmp array back to values
        values[i] = tmp[i-first];
}
```

Merge sort: runtime analysis

- Let's start with a run-time analysis of merge
- Let's use M as the size of the final list
  - The merging requires M (or fewer) comparisons +copies
  - Copying from the temp array is M copies
  - So merge is O(M)

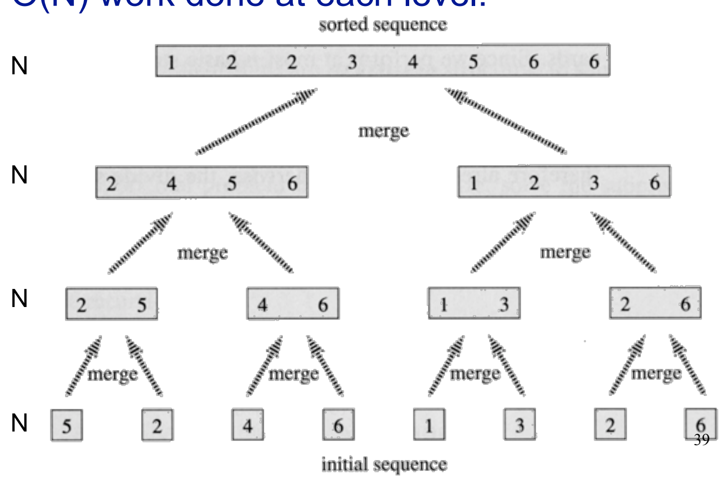
Merge sort: runtime analysis

- The array can be subdivided into halves $\log_2 N$ times (there are $\log_2 N$ levels in the graph)
- At each level in the graph,
  - merge is called on each sub-list
  - The total size of each sub-list added up is N
  - So at each level in the graph, the total execution time is O(N).
- So $\log_2 N$ levels times O(N) at each level:

O(N Log N)

Merge sort
Runtime analysis

- O(N) work done at each level:



Merge sort: runtime analysis

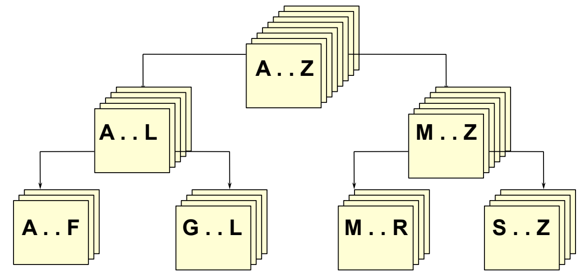
- mergeSort has 2 recursive calls to itself.
- Why does it not have the exponential cost that the Fibonacci algorithm had?

Quicksort

- Another divide and conquer!
- 2 (hopefully) half-sized lists sorted recursively
- the algorithm:
  - If list size is 0 or 1, return. otherwise:
  - partition into two lists:
    - pick one element as the pivot
    - put all elements less than pivot in first half
    - put all elements greater than pivot in second half
  - recursively sort first half and then second half of list.

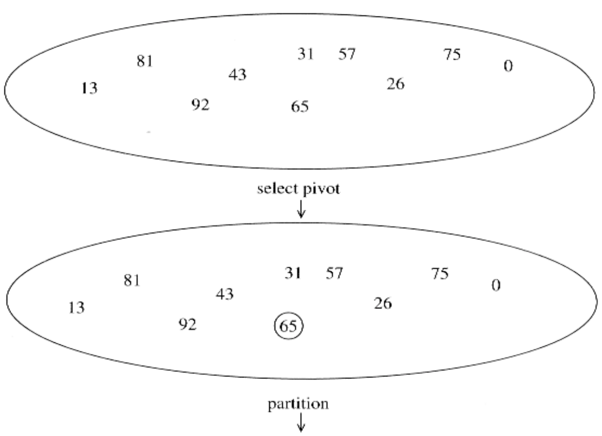
Quicksort

visualization



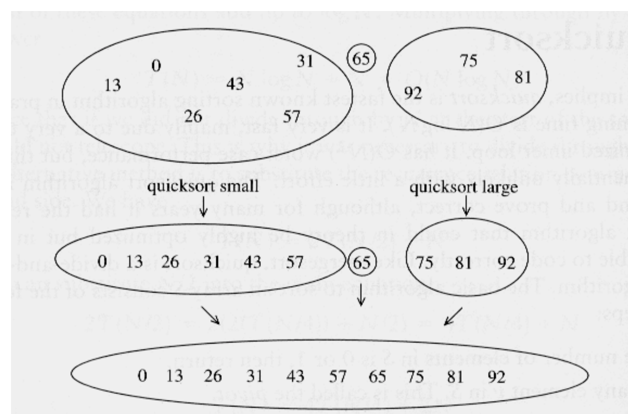
Quicksort

Example



Quicksort

Example cont.



Quicksort: partitioning

- Many algorithms, but must be efficient
- Goal: partition array A [start ... last]
  (start+last are indexes)
- Uses two indexes, i and j, starting from the front and the back of the list
- the algorithm:
  - swap pivot with last element (safekeeping)

5 6 4 6 3 12 19 ➡ 5 6 4 19 3 12 6

pivot i j

Quicksort: partitioning

- the algorithm (continued):
  - move small elements left and larger elements right.
  - let i start at first element, and j start at last-1
  - increment i while A[i] < pivot
  - decrement j while A[j] > pivot

5 6 4 19 3 12 6 ➡ 5 6 4 19 3 12 6

i j i j

Quicksort: partitioning

- the algorithm (continued):
  - When i and j have stopped, and i < j:
  - swap A[i] and A[j]
  - larger element goes to right side, smaller to left
  - maintains: A[x] <= pivot for x<=i and A[x] >= pivot for x >= j

5 6 4 19 3 12 6 $\xrightarrow{\text{swap}}$ 5 3 4 19 6 12 6

i j i j

Quicksort: partitioning

- the algorithm (continued):
  - When i and j have met or crossed (i >= j):
  - swap A[i] and pivot
    - puts pivot in place
    - A[i] >= pivot (i stopped there) so it stays on right side
  - return i (the pivot index)

5 3 4 19 6 12 6

i j

5 3 4 19 6 12 6

j i

5 3 4 6 6 12 19

j i

Quicksort: partitioning

- What if all the elements are bigger than pivot?
  - i never moves, j doesn't stop until it reaches i
  - pivot swapped with A[i], at front of list
  - all elements are to right of pivot
- What if all the elements are smaller than pivot?
  - i never stops until it is at the pivot, j doesn't move
  - pivot swapped with itself, stays at end of list
  - all elements are to left of pivot

Quicksort: partitioning

elements equal to pivot

- What if A[i] or A[j] is equal to the pivot?
- should they stop?
  - if all elements are identical:
    i and j will always stop and swap at every position
  - lots of unnecessary swapping, but pivot ends up in the middle (good).
- if they don't stop:
  - if all elements are identical:
    i never stops until it is at the pivot
  - No swapping, but pivot ends up at end (bad)

Quicksort: code

version 1

```
template<class ItemType>
void quickSort (ItemType values[], int first, int last) {

    if (first < last) {    //at least two elems
        int pivotPoint;

        // partition and get the pivot point (the index)
        pivotPoint = partition(values, first, last);

        quickSort(values, first, pivotPoint - 1);
        quickSort(values, pivotPoint + 1, last);
    }
}

template<class ItemType>
void quickSort (ItemType values[], int size) {
    quickSort(values, 0, size-1);
}
```


Quicksort: code

version 1

```
template<class ItemType>
int partition(ItemType values[], int first, int last) {

    int mid = (first + last) / 2;  //use middle value as pivot

    ItemType pivotValue = values[mid];
    swap(values[last], values[mid]);  //move pivot to end

    int i,j;
    for (i=first, j=last-1; ; ) {
        while (values[i] < pivotValue) {i++;}
        while (j > i && pivotValue < values[j]) {j--;}
        if (i < j) {
            swap(values[i++], values[j--]);
        }
        else
            break;
    }
    swap(values[i], values[last]);   //replace pivot
    return i;
}
```

Quicksort: runtime analysis

- Choice of pivot point dramatically affects running time.
- Best Case
  - Pivot partitions the set into 2 equally sized subsets at each stage of recursion: O(log N) levels
  - Partitioning at each level is O(N)
    - each element is compared to the pivot and maybe moved one time
  - O(N log N)

Quicksort: runtime analysis

- Worst Case
  - Pivot is always the smallest element, partitioning the set into one empty subset, and one of size N-1.
  - Partitioning at each level is N
    - T(N) = T(N-1) + N (time to sort N-1 plus N for partitioning)
    - T(N) = N + N-1 + . . . + 2 + 1 (from unwinding the above)
    - T(N) = N(N+1)/2
  - O(N^2)

Moral of the story: it pays to pick a good pivot point

Quicksort: runtime analysis

- Average Case
  - Assume left side is equally likely to have a size of 0 or 1 or 2 or ... or N elements
  - Partitioning at each level is still N
    - T(N) = average cost of one recursive call, over all subproblem sizes
    - T(N) = (T(0) + T(1) + ... + T(N-1)) / N (divide by N to get avg)
    - Cost for 2 recursive calls and one partitioning:
    - T(N) = N + 2*((T(0) + T(1) + ... + T(N-1)) / N)
    - Not a trivial proof . . . most of it is in the book.
  - O(N log N)

Quicksort: Picking the pivot

- Goal: ensure the worst case doesn't happen.
- Picking a pivot randomly is safe
  - but random number generation can be expensive
- Using the first element:
  - if the input is random, this is ok.
  - if the input is sorted, all elements are in right half worst case = O(N^2)
- Use the median value (the middle value in order):
  - perfectly divides into two even sides
  - but you have to sort the list to find the median.

Quicksort: Picking the pivot

Median of Three method

- Pivot is the median of the first, last, and middle value in the list.
- This is an "estimate" of the real median
  - taking median of more than 3 is not worth the time

Quicksort: Picking the pivot

Median of Three method

- Median-of-Three partitioning:
  - swap the values at first, last and middle so that:

    A[first] = smallest,   A[middle]=median,   A[last] = biggest
  - swap pivot (median) with A[last-1]
  - start with i = first+1 and j=last-2
  - increment i until it encounters an element smaller than pivot
  - decrement j until it encounters an element bigger than pivot
  - if (i<j) swap (A[i], A[j])

Quicksort: Picking the pivot

Median of Three method

2	5	6	4	13	3	12	19	6	A[left] = 2, A[center] = 13, A[right] = 6
2	5	6	4	6	3	12	19	13	Swap A[center] and A[right]
2	5	6	4	6	3	12	19	13	Choose A[center] as pivot
				pivot					
2	5	6	4	19	3	12	6	13	Swap pivot and A[right – 1]
							pivot		

Now we only need to partition A[left + 1, …, right – 2].

Quicksort: Small Arrays

- For very small arrays, quicksort does not perform as well as insertion sort
  - how small depends on many factors, such as the time spent making a recursive call, the compiler, etc
- Do not use quicksort recursively for small arrays
  - Instead, use a sorting algorithm that is efficient for small arrays, such as insertion sort
  - a cutoff between 5 and 20 is good.
  - Note: median of three partitioning requires at least 3 elements anyway

Quicksort: code
version 2

```
template<class ItemType>
void quickSort (ItemType values[], int first, int last) {
    int pivotPoint;
    if (first + CUTOFF <= last) {  // more than CUTOFF elems
        pivotPoint = partition(values, first, last);
        quickSort(values, first, pivotPoint - 1);
        quickSort(values, pivotPoint + 1, last);
    } else {
        insertionSort(values, first,last);
    }
}

template<class ItemType>
void quickSort (ItemType values[], int size) {
    quickSort(values, 0, size-1);
}
```


Quicksort: code
version 2

```
template<class ItemType>
int partition(ItemType values[], int first, int last) {
    //sort first, mid, last
    int mid = (first + last) / 2;
    if (values[mid] < values[first])  swap(values[mid], values[first]);
    if (values[last] < values[first]) swap(values[last], values[first]);
    if (values[last] < values[mid])   swap(values[last], values[mid]);

    ItemType pivotValue = values[mid];  // move pivot to last-1
    swap(values[last-1], values[mid]);

    int i,j;                            // do the partitioning
    for (i=first+1, j=last-2; ; ) {
        while (values[i] < pivotValue) {i++;}
        while (pivotValue < values[j]) {j--;}
        if (i < j)
            swap(values[i++], values[j--]);
        else
            break;
    }
    swap(values[i], values[last-1]);    // put pivot back in place
    return i;
}
```


Quicksort vs MergeSort

- Both run in O(n log n)
- Compared with Quicksort, Mergesort has fewer comparisons but more swapping (copying)
  - (not yet able to verify the following):
  - In Java, an element comparison is expensive but moving elements is cheap. Therefore, Mergesort is used in the standard Java library for generic sorting
  - In C++, copying objects can be expensive while comparing objects often is relatively cheap. Therefore, quicksort is the sorting routine commonly used in C++ libraries